

METHOD AND APPARATUS FOR PROVIDING MULTIPLE QUALITY OF SERVICE LEVELS IN A WIRELESS PACKET DATA SERVICES CONNECTION

BACKGROUND

CROSS REFERENCE OF APPLICATION

This application claims priority from co-pending Provisional Application No. xx/xxx,xxx, filed March 12, 2001, entitled "Method and Apparatus for Providing Multiple Quality of Service Levels in a Wireless Packet Data Services Connection."

Field

[1001] The present invention relates to wireless communications. More particularly, the present invention relates to a novel method and apparatus for providing multiple levels of quality of service in a wireless packet data network between a mobile station and a wireless network.

Background

[1002] The use of code division multiple access (CDMA) modulation techniques is one of several techniques for facilitating communications in which a large number of system users are present. Other multiple access communication system techniques, such as time division multiple access (TDMA), frequency division multiple access (FDMA) and AM modulation schemes such as amplitude companded single sideband (ACSSB) are known in the art. These techniques have been standardized to facilitate interoperability between equipment manufactured by different companies. Code division multiple access communication systems have been standardized in the United States in Telecommunications Industry Association TIA/EIA/IS-95-B, entitled "MOBILE

STATION-BASE STATION COMPATIBILITY STANDARD FOR DUAL-MODE WIDEBAND SPREAD SPECTRUM CELLULAR SYSTEMS", and referred to herein as IS-95. In addition, a new standard for CDMA communication systems has been proposed in the United States in Telecommunications Industry Association (TIA), entitled "Upper Layer (Layer 3) Signaling Standard for cdma2000 Spread Spectrum Systems, Release A — Addendum 1", dated October 27, 2000, and referred to herein as "cdma2000."

[1003] The International Telecommunications Union recently requested the submission of proposed methods for providing high rate data and high-quality speech services over wireless communication channels. A first of these proposals was issued by the Telecommunications Industry Association, entitled "The IS-2000 ITU-R RTT Candidate Submission." A second of these proposals was issued by the European Telecommunications Standards Institute (ETSI), entitled "The ETSI UMTS Terrestrial Radio Access (UTRA) ITU-R RTT Candidate Submission", also known as "wideband CDMA" and hereinafter referred to as "W-CDMA." A third proposal was submitted by U.S. TG 8/1 entitled "The UWC-136 Candidate Submission", hereinafter referred to as "EDGE." The contents of these submissions is public record and is well known in the art.

[1004] IS-95 was originally optimized for transmission of variable-rate voice frames. Subsequent standards have built on the standard to support a variety of additional non-voice services including packet data services. One such set of packet data services was standardized in the United States in Telecommunications Industry Association TIA/EIA/IS-707-A, entitled "Data Service Options for Spread Spectrum Systems", incorporated by reference herein, and hereafter referred to as "IS-707."

[1005] IS-707 describes techniques used to provide support for sending Internet Protocol (IP) packets through an IS-95 wireless network. Packets are encapsulated into a featureless byte stream using a protocol called Point-to-Point Protocol (PPP). Using PPP, IP packets can be transported over a wireless network in segments of arbitrary size. The wireless network maintains PPP state

information for the duration of the PPP session, or as long additional bytes may be sent in the continuous byte stream between the PPP end points.

[1006] Such a continuous byte stream is subsequently encapsulated into a series of IS-95 frames using a protocol called Radio Link Protocol (RLP). RLP includes an error control protocol that uses negative acknowledgments (NAKs) by which the receiver prompts the sender to retransmit lost RLP frames. Because the RLP error control protocol uses retransmissions, RLP data transmission generally exhibits a variable transmission delay from sender to receiver. A modified form of RLP called Synchronous RLP (SRLP), in which no NAKs and no retransmissions are sent by sender or receiver, is well known in the art. The frame error rate in SRLP is greater than that of RLP, but the transmission delay is kept to a minimal constant.

[1007] A remote network node such as a personal or laptop computer (PC) connected to a packet-data-capable wireless mobile station (MS) may access the Internet through a wireless network in accordance with the IS-707 standard. Alternatively, the remote network node such as a web browser may be built-in to the MS, making the PC optional. An MS may be any of a number of types of devices including, but not limited to PC card, personal data assistant (PDA), external or internal modem, or wireless phone or terminal. The MS sends data through the wireless network, where it is processed by a packet data serving node (PDSN). The PPP state for a connection between an MS and the wireless network is typically maintained within the PDSN. The PDSN is connected to an IP network such as the Internet, and transports data between the wireless network and other entities and agents connected to the IP network. In this way, the MS can send and receive data to another entity on the IP network through the wireless data connection. The target entity on the IP network is also called a correspondent node. The interaction between a MS and the PDSN have been standardized in EIA/TIA/IS-835, entitled "Wireless IP Network Standard," dated June, 2000, and referred to herein as "IS-835." One skilled in the art will recognize that, in some networks, the PDSN is replaced with an Interworking Function (IWF).

[1008] In order to provide more complex wireless network services, there is an increasing desire and need to provide different types of services simultaneously through a single wireless device. Examples include simultaneous voice and packet data services. Examples also include multiple types packet data services, such as simultaneous web browsing and video conferencing. At the same time, technological advances are increasing the bandwidth available through a single wireless channel between a wireless device and the wireless network.

[1009] However, modern networks are not yet capable of supporting simultaneous packet data services having substantially different grades of service. For example, delay sensitive applications like video conferencing and voice over IP are optimally sent without RLP retransmissions in order to reduce the magnitude and variability of packet delay through the network. On the other hand, applications such as FTP, e-mail, and web browsing are less delay-sensitive, so are optimally sent using RLP retransmissions. Current wireless standards adequately support a wireless application that requires any one of several grades of service, but not multiple applications in a single MS, where the applications require different grades of service. There is therefore a need in the art for a way of supporting multiple applications in a single MS, where the multiple applications use different grades of service.

SUMMARY

[1010] Embodiments disclosed herein address the above stated needs by enabling a mobile station (MS) and a Radio Access Network (RAN) to establish a connection that supports multiple grades of services with a single IP address assigned to the MS. The embodiments described herein enable a data sender to use a single IP address for multiple packet data applications. The packets generated by each of the multiple packet data applications are provided to a single Point-to-Point Protocol (PPP) stack and a single High-level Data Link Control (HDLC) framing layer to convert data packets into byte streams suitable for transmission through Radio Link Protocol (RLP) connections. Each of the resultant

multiple byte streams is then provided to one of multiple RLP connections having different retransmission and delay properties. The RLP connection selected for sending data from each application is based on the grade of service most appropriate for the application.

[1011] The receiver receives the data on the multiple RLP connections and reassembles the byte streams into frames. The receiver may utilize multiple HDLC framing layers, with one HDLC framing layer corresponding to one RLP connection. Alternatively, the receiver may utilize a single HDLC framing layer and multiple simple "deframer" layers. Each deframer layer corresponds to a single RLP connection, and searches for flag characters that delimit HDLC frames in each RLP byte stream. The deframer layer does not remove HDLC escape codes, but rather provides HDLC stream data to the single HDLC layer as a complete, contiguous HDLC frame.

[1012] The word "exemplary" is used throughout this application to mean "serving as an example, instance, or illustration." Any embodiment described as an "exemplary embodiment" is not to be construed as necessarily preferred or advantageous over other embodiments described herein.

BRIEF DESCRIPTION OF THE DRAWINGS

[1013] FIG. 1 shows an arrangement of protocol layers according to an exemplary embodiment;

[1014] FIG. 2 shows an arrangement of protocol layers according to an alternative embodiment;

[1015] FIG. 3 is a diagram of an exemplary mobile station (MS) apparatus; and

[1016] FIG. 4 is a diagram of an exemplary wireless network apparatus;

[1017] FIG. 5 is a flowchart of an exemplary method of sending packets through multiple RLP connections having different grades of service; and

[1018] FIG. 6 is a flowchart of an exemplary method of receiving packets through multiple RLP connections having different grades of service.

DETAILED DESCRIPTION

[1019] [OPTIONAL DEFINITIONS SECTION]

[1020] Multiple applications using different grades of service could be supported on a single wireless device using a separate Point-to-Point Protocol (PPP) stack for each application. This approach has several disadvantages. Supporting multiple PPP instances for a single mobile station (MS) would needlessly consume large amounts of data memory in both the MS and the Packet Data Serving Node (PDSN).

[1021] In addition, if a Radio Link Protocol (RLP) session were established for use by an application that required low latency, the RLP should be configured to operate without retransmissions. While this would result in the low latency that is best for overlying application, the Link Control Protocol (LCP) and other configuration protocols that are required to establish the PPP link would have to proceed without error control. The resulting increase in frame error rate could cause delays or even failure of PPP configuration before any application packets could be sent.

[1022] The embodiments discussed below overcome these disadvantages by utilizing a single PPP instance for multiple RLP instances between a MS and the wireless network. **FIG. 1** shows an arrangement of protocol layers between a sender and a receiver of packet data using different concurrent grades of service. In an exemplary embodiment, the sender maintains two Radio Link Protocol (RLP) layers (**106** and **108**), one High-level Data Link Control (HDLC) layer **104** and one Point-to-Point Protocol (PPP) layer **102**. Each of the RLP layer instances uses different grades of service (**106** and **108**). For example, if RLP_{1s} **106** is configured to retransmit frames in response to NAK frames received from the receiver, RLP_{2s} **108** is configured for no retransmissions. In other words, RLP_{1s} **106** provides higher reliability through use of an error control protocol, while RLP_{2s} **108** provides unreliable transport with fixed, minimal transmission delay. The grade of service that characterizes such an RLP_{1s} **106** is referred to herein as “reliable” for short. Similarly, the grade of service that characterizes such an RLP_{2s} **108** is referred to

herein as “low latency.” Though exemplary embodiments are described herein as using just two grades of service, implementations that use a larger number of different grades of service are also anticipated and are to be considered within the scope of the described embodiments. For example, a sender and receiver could each additionally utilize a third RLP layer that provides an intermediate grade of service having a degree of reliability that lies between “reliable” and “low latency.”

[1023] In an exemplary embodiment, the receiver also maintains two receive RLP instances (**116** and **118**) corresponding to the same grades of service as the RLP instances in the sender (**106** and **108**). For example, if RLP_{1S} **106** provides reliable grade of service, then RLP_{1R} **116** is configured for reliable grade of service. Thus, when the RLP_{1R} **116** layer detects a break in the sequence numbers of the received RLP frames, then RLP_{1R} **116** responds by sending a NAK frame to prompt retransmission. Upon receiving an RLP NAK frame, RLP_{1S} **106** retransmits the requested frame from its retransmission buffer. On the other hand, if RLP_{2S} **108** is configured for low latency grade of service, then RLP_{2R} **118** will not send a NAK frame regardless of breaks in frame sequence numbers. Indeed, RLP_{2S} **108** and RLP_{2R} **118** may omit frame sequence numbers entirely from the transmitted RLP frames to make more room for data payload. In addition, RLP_{2S} **108** need not maintain a retransmission buffer of previously sent frames, thus conserving memory in the sender. Also, RLP_{2R} **118** need not maintain a resequencing buffer, thus conserving memory in the receiver.

[1024] The PPP_S layer **102** in the sender encapsulates IP packets within PPP frames. In an exemplary embodiment, the PPP_S layer **102** increases packet throughput by performing IP header compression such as the well-known Van-Jacobsen (VJ) header compression. VJ header compression can result in the loss of certain header information that might otherwise be useful in multiplexing PPP packets between the multiple RLP layers (**106** and **108**). In an exemplary embodiment, the PPP_S layer **102** provides whole PPP packets to the HDLC_S layer **104**, and also provides information that can be used to determine which RLP layer to send the framed data through. In an exemplary embodiment, the PPP_S layer **102** provides a grade of service identifier or an RLP instance identifier with each

PPP packet provided to the HDLC_S layer **104**. The HDLC_S **104** layer adds flag characters between the PPP packets and adds a cyclical redundancy checksum (CRC) to each PPP packet received from the PPP_S layer **102**. The HDLC_S layer **104** further performs HDLC escaping to ensure that no flag or HDLC control characters appear within the data of a single frame. The HDLC_S layer **104** typically performs HDLC escaping by replacing each flag or control character with an escape sequence having at least two characters.

[1025] The receiver in **FIG. 1** is shown with a separate HDLC layer (**112** and **114**) for each RLP instance (**116** and **118**). The bytes received in the RLP frames by each RLP instance (**116** and **118**) are presented to the corresponding HDLC layer instances (**112** and **114**). Each HDLC layer instance (**112** and **114**) locates escape sequences in its respective input data stream and converts each escape sequence back to the original data within the transmitted frames. The HDLC layer instances (**112** and **114**) also perform checks of the CRC's received in the frames to determine whether the frames were received with communication errors. Frames having incorrect CRCs are silently discarded, and frames having correct CRC's are forwarded to the next protocol layer up (PPP_R) **110**.

[1026] **FIG. 2** shows an alternate arrangement of protocol layers. The arrangement of protocol layers in the sender in **FIG. 2** is identical to that of the sender in **FIG. 1**. In the receiver, however, a single HDLC_R layer **212** is used instead of one for each RLP instance. Deframer layers (**214** and **220**) are inserted between the RLP layers (**218** and **216**) and the HDLC_R layer **212**. The purpose of the deframers (**214** and **220**) is to ensure that only whole HDLC frames are delivered to the HDLC_R layer **212**. Delivering only whole HDLC frames makes it unnecessary for the HDLC_R layer **212** to differentiate between, or reassemble, data from multiple HDLC frames. The HDLC_R layer **212** removes escape sequences and checks the CRC for an entire frame. If the CRC is deemed correct, then the HDLC_R layer **212** forwards the complete PPP frame to the PPP_R layer **210**. If the CRC is incorrect, the HDLC_R layer **212** silently discards the erroneous frame data.

[1027] One benefit of using deframer layers (**214** and **220**) is that it enables the receiver to support multiple instances of RLP (**218** and **216**) without any changes

to the implementation of the HDLC_R layer **212**. The HDLC_R layer **212** need not even be aware that the received bytes have been received through two different RLP connections. This implementation independence is especially important in network implementations where the HDLC_R protocol layer **212** resides in a different physical device than the RLP protocol layers. For example, the HDLC_R layer might exist within a standard packet router, and the RLP layers might exist within a Packet Control Function (PCF) within a Radio Access Network (RAN) of a wireless network. The use of the deframer layers makes it possible to support multiple RLP layers and grades of service without changing the software of the standard packet router.

[1028] FIG. 3 shows an exemplary mobile station (MS) that supports the multiple grades of service as discussed above. A control processor **302** establishes a wireless connection through a wireless modem **304**, transmitter **306**, and antenna **308** as shown. In an exemplary embodiment, the wireless modem **304** and transmitter **306** operate in accordance with the cdma2000 specification. Alternatively, the wireless modem **304** and transmitter **306** could operate in accordance with some other wireless standard such as IS-95, W-CDMA, or EDGE.

[1029] The control processor **302** is connected to a memory **310** having code or instructions directing the control processor **302** to establish and utilize the protocol layers shown in FIGS. 1-2. The memory **310** may include RAM memory, flash memory, ROM memory, EPROM memory, EEPROM memory, registers, hard disk, a removable disk, a CD-ROM, or any other form of storage medium or computer readable media known in the art.

[1030] In an exemplary embodiment, the control processor **302** uses some of the memory **310** as memory buffers (**312** and **314**) necessary for operation of the multiple RLP layers. For example, if the RLP₁ buffer **312** corresponds to a reliable RLP connection, it will include a retransmission buffer for RLP data being sent and a resequencing buffer for RLP data being received. If the RLP₂ buffer **314** corresponds to a low latency RLP connection, then the RLP₂ buffer **314** need not have either a retransmission buffer or a resequencing buffer. Because these two buffers are not needed, the RLP₂ buffer **314** occupies a smaller amount of memory

than the RLP₁ buffer 312. Though depicted as disjoint, the buffers (312 and 314) may also overlap if some data structures are shared between the multiple RLP implementations.

[1031] FIG. 4 shows an exemplary wireless communication network having a connection with a packet network such as the Internet 416. The wireless communication network includes a RAN 412 and a PDSN 414. The RAN 412 further includes a selector 402 that is connected to one or more wireless base stations (not shown). The selector 402 in the RAN 412 is generally a subsystem of a base station controller (BSC), which is not shown. All wireless data sent to or received from the MS is routed through the selector. In addition to the selector 402, the RAN 412 also includes a Packet Control Function (PCF) 404. For packet data service options, the selector sends packet data received from the MS through the PCF 404, which further includes a control processor 460 and memory 418.

[1032] The memory 418 contains code or instructions directing the control processor 460 to establish and utilize the protocol layers shown in FIGS. 1-2. The memory 418 may include RAM memory, flash memory, ROM memory, EPROM memory, EEPROM memory, registers, hard disk, a removable disk, a CD-ROM, or any other form of storage medium or computer readable media known in the art.

[1033] In an exemplary embodiment, the control processor 406 establishes multiple buffer areas (408 and 410) within the memory 418 for the various RLP connections established with multiple mobile stations. In an exemplary embodiment, one pool of RLP₁ buffers 408 includes retransmission and resequencing buffers for use in reliable RLP instances. Another pool of RLP₂ buffers 410 is used for low latency RLP instances, and therefore does not include retransmission and resequencing buffers. The control processor 406 may allocate more than one RLP instance to a single MS. For example, one RLP₁ buffer and one RLP₂ buffer may be allocated to a single MS that is running a combination of delay-sensitive and non-delay-sensitive applications.

[1034] The control processor 406 is also connected to a PDSN 414. In an exemplary embodiment, when the MS sends an IP packet to the packet network 416, the control processor 406 receives the RLP frames from the selector 402 and

uses the associated RLP buffer (408 or 410) to extract a stream of bytes from the RLP frames. The bytes are then sent from the control processor 406 to the PDSN 414, which extracts complete IP packets (those having correct CRC values) from the byte stream in accordance with the HDLC protocol. The PDSN 414 then forwards the resultant IP packets to the packet network 416. If the PDSN 414 maintains a single HDLC connection for multiple RLP connections to a single MS, then the control processor 406 performs deframing before sending the bytes from the RLP frames to the PDSN 414. The result of deframing is that whole HDLC frames are forwarded by the control processor 406 to the PDSN 414. In other words, the control processor 406 ensures that no data from an HDLC frame received on one RLP link will be intermingled with the data from an HDLC frame received on another RLP link. Deframing allows better resource usage in addition to enabling the use of existing PDSNs that cannot allocate more than a single PPP/HDLC to an IP address.

[1035] When the packet network 416 sends a packet to the MS, the packets are first received at the PDSN 414. In an exemplary embodiment, the PDSN 414 encapsulates the IP datagrams addressed to the MS into PPP packets and uses HDLC framing to convert the resulting PPP packets into a stream of bytes. In an exemplary embodiment, the PDSN 414 assigns a single HDLC instance to a single MS, and uses that HDLC instance to perform HDLC framing of any IP packet addressed to that MS. In an alternate embodiment, the PDSN 414 may have multiple HDLC instances assigned to a single MS, such that each HDLC instance corresponds to a single RLP connection within the MS.

[1036] The connections between the PDSN 414 and the network 416, between the PDSN 414 and the control processor 406, and between the control processor 406 and the selector 402, may use any of a variety of interfaces including ethernet, T1, ATM, or other fiber, wired or wireless interface. In exemplary embodiment, the connection between the control processor 406 and the memory 418 will generally be a direct hardware connection such as a memory bus, but may also be one of the other types of connection discussed above.

[1037] **FIG. 5** is a flowchart of an exemplary method of sending packets through multiple RLP connections having different grades of service. In an exemplary embodiment, a control processor of the sending apparatus (**302** of **FIG. 3** or **406** of **FIG. 4**) utilizes the method characterized by **FIG. 5**. At step **502**, the sender encapsulates the IP packet to be sent into a PPP packet. In an exemplary embodiment, IP header compression such as Van-Jacobson (VJ) header compression is also performed at step **502**. At step **504**, the sender then converts the PPP packet into a byte stream according to the HDLC protocol. Specifically, each PPP packet is converted into an HDLC frame. One or more flag characters are inserted between HDLC frames in the byte stream, and flag and control characters that appear within each frame are replaced with escape sequences. Probably the most common examples of HDLC escaping are replacing the flag sequence octet 0x7e (hexadecimal) with the two octets 0x7d 0x5e (hexadecimal) and replacing octet 0x7d (hexadecimal) with the two octets 0x7d 0x5d (hexadecimal). Also at step **504**, a CRC is computed for each frame and inserted at the end of the frame (before the flag character that signals the end of the frame). At step **506**, the sender determines which of a set of available grades of service should be used to send the data for the frame based on the type of the packet. IP packets being sent using non-delay-sensitive applications such as FTP or TCP are sent at step **508** using reliable RLP (with retransmission and resequencing). Also, any packets that are not IP packets, but still non-delay-sensitive (such as IPCP or LCP packets) are sent at step **508** using reliable RLP. Delay-sensitive types of packets, such as Real-Time Protocol (RTP) packets used for video conferencing services, are sent at step **510** using low latency RLP. As discussed above, low latency RLP does not send or request retransmissions of RLP frames lost to communication errors. Though two grades of service are shown in the exemplary embodiment of **FIG. 5**, one skilled the art will recognize other systems may utilize more than two different grades of service without departing from the scope of the described embodiments. For example, at step **506**, the sender may elect to send some types of packets through an RLP connection having an intermediate level of reliability.

[1038] FIG. 6 is a flowchart of an exemplary method of receiving packets through multiple RLP connections having different grades of service. In an exemplary embodiment, a control processor of the receiving apparatus (302 of FIG. 3 or 406 of FIG. 4) utilizes the method characterized by FIG. 6. At step 602, the receiver processes received RLP frames received through one or more RLP connections. In an exemplary embodiment as described above, the RLP frames are received through two types of RLP connections, low latency and reliable.

[1039] As described in the aforementioned IS-707, RLP frames received through a reliable RLP connection have sequence numbers that the receiver uses to resequence the frames and request retransmission of lost frames. For example, if an RLP frame having the sequence number "7" is lost to communication errors, the receiver sends a NAK frame to request retransmission of that frame. When the retransmitted frame is received, the data carried in that frame is used to complete the stream of data bytes before providing any subsequent data bytes to the HDLC layer. As a result, the stream of data bytes extracted from the RLP frames of a reliable RLP connection will generally not have gaps compared to what was transmitted by the sender. The cost of preventing gaps in the data is variable latency.

[1040] In contrast, when an RLP frame is lost to communication error on a low latency RLP link, no retransmission is requested or sent. Any data bytes carried in such a lost RLP frame is omitted from the stream of data bytes presented to the receiver's HDLC layer. In other words, the loss of an RLP frame on a low latency RLP link invariably causes a gap in the receiver's stream of data bytes compared to what was transmitted by the sender. However, low latency RLP protocol has latency that is both fixed and small, making it highly suitable for sending delay-sensitive types of packets, such as RTP packets.

[1041] In an exemplary embodiment described in FIG. 2, the receiver uses deframers (214 and 220 in FIG. 2) received through multiple RLP connections (116 and 118 in FIG. 2) in order to provide whole HDLC frames of data to a single HDLC protocol layer (212 in FIG. 2). In FIG. 6, this deframing is performed at step 604. At step 606, the HDLC protocol layer (212 in FIG. 2) removes HDLC escape

sequences that were inserted by the sender and checks the CRC of each HDLC frame. At step **606**, any HDLC frame bearing an incorrect CRC is silently discarded by the receiver. The resulting PPP frames are then provided by the HDLC protocol layer to the PPP layer. At step **608**, the PPP layer decapsulates the received packets, removing the PPP header and any other changes made by the sender. Also at step **608**, if the sender compressed the IP header of the received packet (for example, using VJ header compression), then the IP header is expanded to its original size and contents. The decapsulated packets are then routed at step **610**. Though the embodiments described above discuss primarily encapsulating IP packets, PPP and HDLC can also be used to send packets for other protocols such as IPX or LCP.

[1042] In an exemplary embodiment using deframers (**214** and **220** in **FIG. 2**), steps **602** and **604** are performed by the control processor (**406** of **FIG. 4**) within the RAN (**412** of **FIG. 4**), and steps **606**, **608**, and **610** are performed by the PDSN (**414** of **FIG. 4**). In alternate embodiment such as that shown in **FIG. 1**, the PDSN (**414** of **FIG. 4**) allocates multiple HDLC layers (**112** and **114** of **FIG. 1**) to a single MS. In this embodiment, there is no deframing performed by the receiver, and step **604** is omitted. At step **602**, each RLP layer (**116** and **118** of **FIG. 1**) provides data extracted from received RLP frames directly to its corresponding HDLC layer (**112** and **114** of **FIG. 1**, respectively).

[1043] Thus, described herein is a method and apparatus for providing multiple quality of service levels in a wireless packet data services connection. Those of skill in the art would understand that information and signals may be represented using any of a variety of different technologies and techniques. For example, data, instructions, commands, information, signals, bits, symbols, and chips that may be referenced throughout the above description may be represented by voltages, currents, electromagnetic waves, magnetic fields or particles, optical fields or particles, or any combination thereof. One skilled in the art will also recognize that the PDSN in the above-described embodiments could also be replaced with an Interworking Function (IWF) without departing from the scope of the described embodiments.

[1044] Those of skill would further appreciate that the various illustrative logical blocks, modules, circuits, and algorithm steps described in connection with the embodiments disclosed herein may be implemented as electronic hardware, computer software, or combinations of both. To clearly illustrate this interchangeability of hardware and software, various illustrative components, blocks, modules, circuits, and steps have been described above generally in terms of their functionality. Whether such functionality is implemented as hardware or software depends upon the particular application and design constraints imposed on the overall system. Skilled artisans may implement the described functionality in varying ways for each particular application, but such implementation decisions should not be interpreted as causing a departure from the scope of the present invention.

[1045] The various illustrative logical blocks, modules, and circuits described in connection with the embodiments disclosed herein may be implemented or performed with a general purpose processor, a digital signal processor (DSP), an application specific integrated circuit (ASIC), a field programmable gate array (FPGA) or other programmable logic device, discrete gate or transistor logic, discrete hardware components, or any combination thereof designed to perform the functions described herein. A general purpose processor may be a microprocessor, but in the alternative, the processor may be any conventional processor, controller, microcontroller, or state machine. A processor may also be implemented as a combination of computing devices, e.g., a combination of a DSP and a microprocessor, a plurality of microprocessors, one or more microprocessors in conjunction with a DSP core, or any other such configuration.

[1046] The steps of a method or algorithm described in connection with the embodiments disclosed herein may be embodied directly in hardware, in a software module executed by a processor, or in a combination of the two. A software module may reside in RAM memory, flash memory, ROM memory, EPROM memory, EEPROM memory, registers, hard disk, a removable disk, a CD-ROM, or any other form of storage medium or computer readable media known in the art. An exemplary storage medium is coupled to the processor such the

processor can read information from, and write information to, the storage medium. In the alternative, the storage medium may be integral to the processor. The processor and the storage medium may reside in an ASIC. The ASIC may reside in a mobile station. In the alternative, the processor and the storage medium may reside as discrete components in a mobile station.

[1047] The previous description of the disclosed embodiments is provided to enable any person skilled in the art to make or use the present invention. Various modifications to these embodiments will be readily apparent to those skilled in the art, and the generic principles defined herein may be applied to other embodiments without departing from the spirit or scope of the invention. Thus, the present invention is not intended to be limited to the embodiments shown herein but is to be accorded the widest scope consistent with the principles and novel features disclosed herein.

[1048] WHAT IS CLAIMED IS: